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Please do not open this question set before you are allowed to do so.**

4810-1184 Algorithms for Information Security and Privacy (Autumn 2018)

Midterm Problem 1

Suppose that we have a table T with the weight information of all persons in a city. We want to use the Laplacian mechanism to publish the average weight information $AVG(T)$. However, for now, let suppose that $f(T) = AVG(T)/2$.

Question 1.1: Calculate the value of $GS(f)$.

Question 1.2: What is the distribution that the noise we add to $f(T)$ should be drawn from?

Question 1.3: Discuss why your noise in Question 1.2 should be smaller than when $f(T) = AVG(T)$.

Question 1.4: Discuss why, even the noise is smaller, scaling the value of f will not help users to have a clearer information.

From next question, we will use the exponential mechanism to privately publish the average weight $AVG(T)$. Suppose that a publication of the mechanism p must be a member of the set $\{40,50,60,70,80,90\}$.

Question 1.5: Design a utility function for each p .

Question 1.6: Calculate the value of $\Delta Utility$.

Question 1.7: Discuss why, when using your utility function, the exponential mechanism will work well.

Question 1.8: Discuss why publishing $AVG(T)/2$ instead of $AVG(T)$ will also not work in the exponential mechanism with your utility function. You can assume that, when trying to publish $AVG(T)/2$, the possible publication can be $\{20,25,30,35,40,45\}$.

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Midterm Problem 2

Suppose that our database T has 2 information, chosen party and ages, of all users. Also, for simplicity, let assume that there is only 2 parties in our country, Melon-pan and Gyu-don. We know in advance that there will be two queries to our database, which are:

- 1) How many persons choose “Melon-pan” and has age below 30?
- 2) What is the percentage of persons who is below 40 and choose “Melon-pan” party?

Question 2.1: Discuss why reducing the table T to T' with the following information:

- 1) chosen party
- 2) if the person ages below 30 (boolean variable)
- 3) if the person ages below 40 (boolean variable)

will help reducing the computation time of the smallDB algorithm.

Question 2.2: Design a function f_1 for Query 1. Your function must follow the assumption on the function for smallDB algorithm discussed in the class.

Question 2.3: Design a function f_2 for Query 2. Your function must follow the assumption on the function for smallDB algorithm discussed in the class.

From next question, suppose that $\alpha = 0.3$.

Question 2.4: Calculate the number of records in the published small database, when $\log 2 = 0.7$.

Question 2.5: Calculate the number of possible small tables.

From next question, assume that, the histogram of the table T' is as follows:

- Number of records with “Melon-pan” and “age below 30” is 500000
- Number of records with “Melon-pan” and “age over 30” is 0
- Number of records with “Gyu-don” and “age below 40” is 0
- Number of records with “Gyu-don” and “age over 40” is 500000

Question 2.6: Calculate the number of possible small table U with $\max_i [f_i(T') - f_i(U)] = 0$.

Question 2.7: Calculate the number of possible small table U with $\max_i [f_i(T') - f_i(U)] = 1/8$.

Question 2.8: Based on your answer in Questions 2.6 and 2.7, try to a fast algorithm for smallDB algorithm. By the algorithm, can you reduce the value of α to a smaller value?

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Midterm Problem 3

Suppose that we want to publish the percentage of persons choosing Melon-pan party, i.e.

$$f(T) = \frac{\#Melon - pan}{|T|}.$$

We know that $f(T)$ must be between 0 and 1.

Question 3.1: Discuss why Laplacian mechanism can give a number smaller than 0 or larger than 1.

Question 3.2: Discuss why there would not be a privacy problem although the number is not between 0 and 1.

From next question, let consider the case that it is very unlikely to have $f(T) + noise$ smaller than 0 or larger than 1. The probability that we have such the case it negligible.

Suppose that there is an attacker who want to know $f(T)$. The attacker keeps asking from the value $f(T) + noise$ for n times. Assume that he has got X_1, \dots, X_n for the queries.

Question 3.3: Discuss why the expected value of X_i is $f(T)$.

Let S be an average of X_i . The attacker will predict that $f(T)$ is S .

Question 3.4: Discuss why the expected value of S is $f(T)$.

Question 3.5: Discuss why the probability that S is different from $f(T)$ by more than α is no more than $2 \cdot \exp(-2n\alpha^2)$.

Question 3.6: Calculate the probability that the attacker's prediction has error smaller than $\sqrt{\frac{\ln 10}{200}} \approx 0.11$ when $n = 100$.

Question 3.7: Discuss why your answer in Question 3.6 would imply a privacy problem for Laplacian mechanism. Discuss why we have already included this problem in the composition theorem.