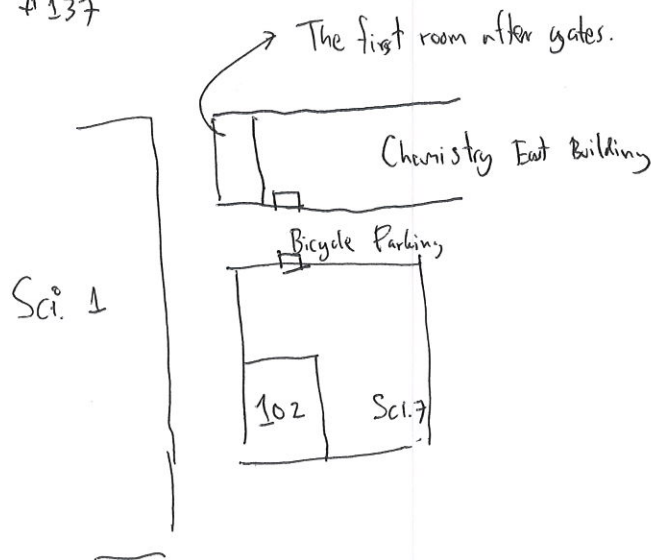


Vorapong Suppatitpaisarn (International Center for IST)

* I sometimes announce about exchange programs with foreign universities at this class.

Office hour : Thursday 2:30 - 4:00 pm.

Office : Chemistry Building #137



There is no room number pasted in front of my room!!

Schedule

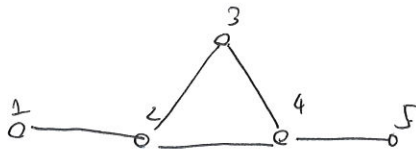
- | | |
|---------------------|--|
| 10/1 | Introduction, Quick catch up on our course on last semester. |
| 10/8 | No class (public holiday). |
| 10/15 | Social Sensor Networks: Graph Clustering |
| 10/22, 10/29 | Social Networks: Page Rank and SimRank |
| 11/5, 11/8
(Thu) | Social Networks: Influence Maximization. |
| 11/12 | Midterm Examination (30% of credits) |
| 12/19 | Sensor Networks: Target Coverage Problem |
| 11/26, 12/3 | Sensor Networks: Maximum Lifetime Coverage Problem |
| 12/10 | Sensor Networks: Localization Problem |
| 12/18 | Guest Lecture (TBP) |
| 12/24, 12/31 | No class. |
| 1/7 | Final Examination (70% of credits) |

Please contact me before 10/15 if you can not join the midterm and final exam!

Graph V : Set of points (nodes)

E : Set of lines between points (edges or links)

Example



$V = \{1, 2, 3, 4, 5\}$

$E = \{\{1, 2\}, \{2, 3\}, \{2, 4\}, \{3, 4\}, \{4, 5\}\}$

Paths

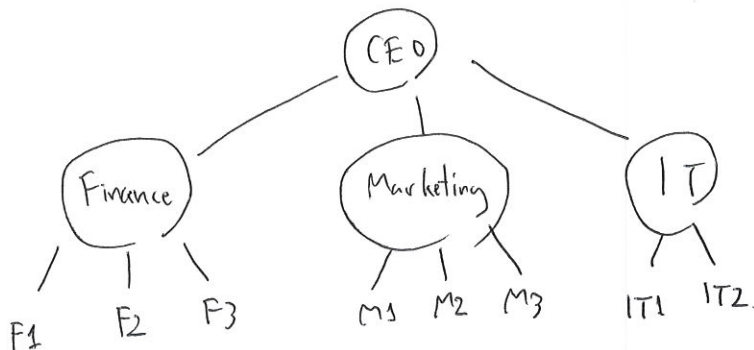
1 wants to communicate with 5.

1 can send a message to 2 asking him to ask 4 to send to 5.

paths $\left(\begin{array}{l} 1 \rightarrow 2 \rightarrow 4 \rightarrow 5 \\ 1 \rightarrow 2 \rightarrow 3 \rightarrow 4 \rightarrow 5 \end{array} \right. \rightsquigarrow \text{paths with smallest number of forwarding} = \text{shortest path}$

Trees A graph that there is only one path between all two points.

Example organization chart



Social Networks: Networks of persons

V : points \rightarrow persons

E : lines \rightarrow lines between 2 persons that are friends.

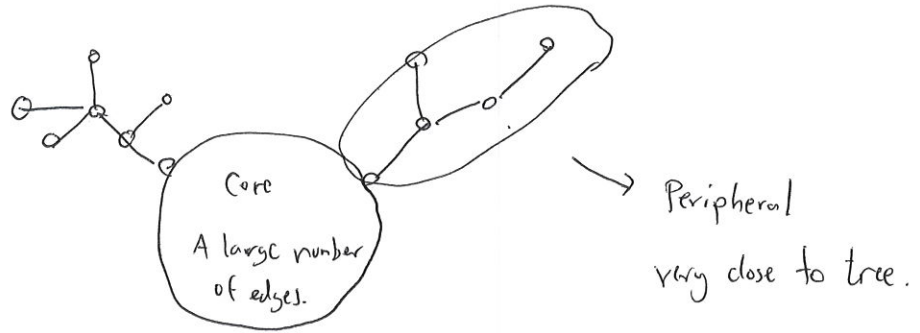
What makes social networks different?

1. Small-world Properties.

I am a friend of a friend of a friend of a friend of a friend of a friend of everyone.

Shortest paths between 2 persons are not longer than 6 forwardings.

2. Core-peripheral Property

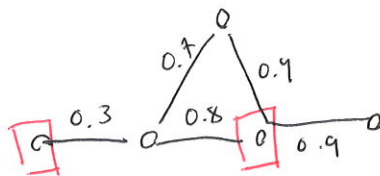


Use 2 properties to speed up computations.

Influence Maximization.

- We want to advertise an information in a social network
- We have a budget to tell at most k (2-3) persons in the network, and hope that they will forward the information to their friends as much as possible.

Example

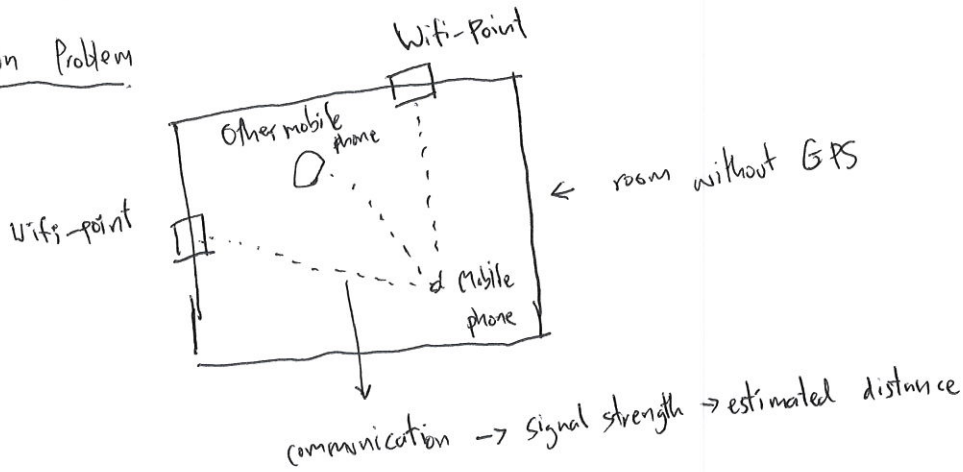


$k=2$

Sensor Networks

- Dense networks, Each node has small energy
- Management has to be seriously done.

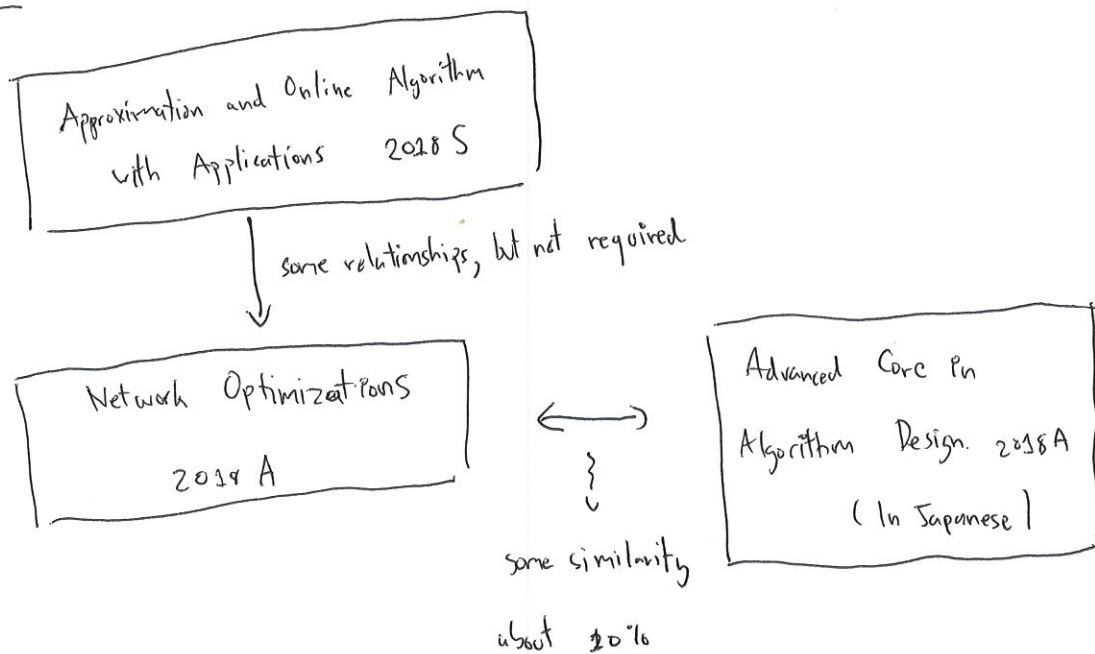
Localization Problem



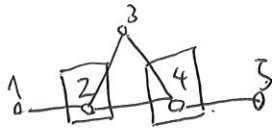
Technical Concepts

- k-way clustering
- Submodular function optimization.
- Garg-Komman. optimization method
- Semi-Definite Programming (SDP)

Relationship with other courses



Domination Set



- We want everyone to know the information
- We select persons, they will forward information to all friends. The friends will not further forward the information
- We will select the smallest number of persons.



Optimization Formulation

Input: Social Network (V, E)

What the output can be?

Output: Set of selected persons $S \subseteq V$



Constraint: For all $v \in V$, either $v \in S$ or there is $u \in S$ such that $\{u, v\} \in E$

Objective Function: Minimize $|S|$ \rightarrow What is the desirable output?



Input: Social Network (V, E)

Output: Suppose that $V = \{1, \dots, n\}$, $x_1, x_2, \dots, x_n \in \{0, 1\}$

~~$x_i = 0$~~ $x_i = 0$ if we do not select i
 $x_i = 1$ if we select

Objective Function: Minimize $\sum_{i=1}^n x_i$

Constraint: For all i , $\sum_{j \in N(i) \cup \{i\}} x_j \geq 1$
 \downarrow
 friends of i

Example of constraints

$$x_2 + x_3 \geq 1$$

\leftarrow 1 must know information

$$x_1 + x_2 + x_3 + x_4 \geq 1$$

$$x_2 + x_3 + x_4 \geq 1$$

$$x_2 + x_3 + x_4 + x_5 \geq 1$$

$$x_4 + x_5 \geq 1$$